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### Abstract

First, the paper presents in short algorithms for simulating some distributions which are used in the paper. Then, some probability distributions are introduced, which are mixtures between distributions of a minim and a maxim of sequences of life data having Lomax (Pareto) or Weibull distributions (called target distributions), mixed up with truncated geometric or Poisson distributions. These distributions could be used to calculate the reliability of multicomponent serial or paralel systems and particularly for simulating such sistems. Apart from composition methods for simulating these distributions, some other simulation algorithms, based on inverse or rejection methods, are presented. **Keywords:** distribution functions

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### 1. INTRODUCTION

In [1,2,3] were introduced some mixtures of life data distributions, deriving from some known distributions, which involve mixtures with Poisson distribution. Here we continue to introduce similar distributions which involve further mixtures with Poisson distributions or geometric distributions. More precisely, we consider a life time L having a target cumulative distribution function  $(cdf) \phi(x)$  and denote  $\varphi(x)$  its probability density function (pdf). Hence  $\varphi(x) = \phi'(x)$ .

If  $L_1, L_2, ..., L_n$  are independent and identically distributed *(iid)* random variables, then, let us consider the *extremum ranking* variables, namely,

$$V = \min_{1 \le i \le n} L_i; \ W = \max_{1 \le i \le n} L_i.$$
(1.1)

Note that the random variables V and W have the following cdf's respectively

$$\phi_V(x) = (1 - (1 - \phi(x))^n; \ \phi_W(x) = (\phi(x))^n \tag{1.2}$$

and the corresponding pdf's are

$$\varphi_V(x) = n\varphi(x)(1-\phi(x))^{n-1}; \ \varphi_W(x) = n\varphi(x)(\phi(x))^{n-1}.$$
 (1.3)

In this paper we consider as target distributions the  $Lomax(a, \theta)$  and the *standard Weibull*( $\nu$ ) distributions. The sample size n is assumed to be a discrete random variable  $N^*$  which takes values  $N^* = 1, 2, ...$  In the following  $N^*$  is either a truncated (on  $[1, \infty)$ ) Poissaon( $\lambda$ ) random variable, or a truncated Geometric(p),  $0 random variable. In the Poisson(<math>\lambda$ ) case, the truncated distribution is

$$P(N^* = k) = \frac{1}{e^{\lambda - 1}} \frac{\lambda^k}{k!}, \ \lambda > 0, \ k = 1, 2, \dots$$
(1.4)

and in the Geometric(p) case, the truncated distribution is

$$P(N^* = k) = pq^{k-1}, \ k = 1, 2..., \ 0 (1.5)$$

The mixed distribution will be obtained as distributions of V or W with random  $n = N^*$  (i.e. truncated  $Poisson(\lambda)$  or Geometric(p)).

#### 1.1 Simulation of discrete distributions involved

• The  $Poisson(\lambda)$  distribution of the discrete random variable N is

$$P(N=k) = \frac{\lambda^k}{k!} e^{-\lambda}, \ k = 0, 1, 2, \dots$$
(1.6)

There are several known ways of simulating this distribution. One convenient way is based on the following property:

If  $n \in \mathcal{N}^+$  and  $p, 0 are selected such as <math>n \to \infty$  (i.e. n is large) and  $\lambda = n.p$  then  $Poisson(\lambda)$  distribution can be approximated with the Binomial Bin(n, p) distribution.

The Bin(n, p) distribution is in the form

$$P(N = k) = C_n^k p^k q^{n-k}, q = 1 - p, k = 0, 1, ..., n$$
(1.7)

As n is large it is also known that

The random variable

$$Z = \frac{N - np}{\sqrt{npq}} \tag{1.8}$$

is approximately distributed as normal N(0,1)

As the normal N(0,1) random variable can be easily simulated (see for instance [4],p.68-69 and & 3.5.4), the random variable N can be simulated as

Generate  $Z \to N(0,1)$ ; Take  $N = \{np + Z\sqrt{npq}\}$ .

(The notation  $\{x\}, \forall x \in R$  means the closest integer from x).

• The Geometric Geom(p) (0 random variable N has the following distribution

$$P(N = k) = pq^k, \ k = 0, 1, \dots, q = 1 - p$$
(1.9)

Note that the *cumulative* geometric distribution function is

$$F(x) = P(N < x) = \sum_{i=1}^{x-1} pq^i = 1 - q^x, \ x - 0, 1, 2, \dots$$
(1.10)

therefore, the simulation of N in this case can be done by the inverse method as follows

Generate U an uniform (0,1) random number;

Take 
$$N = \left[\frac{\log(U)}{\log(q)}\right]$$
.

(The  $[x], \forall x \in R$ , denotes the integer part of x).

The described algorithms simulate discrete random variables N having  $Poisson(\lambda)$  or Geometric Geom(p) distributions. For the purpose of this paper, we need to simulate *truncated* random variables  $N^* = N, N > 0$ . Any of these random variables are easily genereted by the following *rejection* algorithm.

#### repeat

Generate  $N \rightarrow Poisson(\lambda)$  (or Geom(p)); until N > 0; Take  $N^* = N$ .

## 1.2 Simulation of discrete distributions involved

As concerns Weibull distribution it is enough to consider only the standard Weibull $(0,1,\nu)$  distribution having the pdf in the form

$$f(x,\nu) = \begin{cases} 0 & \text{if } x \le 0\\ \nu x^{\nu-1} e^{-x^{\nu}} & \text{if } x > 0. \end{cases}$$
(1.11)

The corresponding cdf is

$$F(x,\nu) = \begin{cases} 0 & \text{if } x \le 0\\ 1 - e^{-x^{\nu}} & \text{if } x > 0, \end{cases}$$
(1.12)

and its *inverse* is the solution of the equation F(x) = U i.e.

$$F^{-1}(x) = (-\log U)^{\frac{1}{\nu}}, \text{ for } 0 < U < 1.$$
 (1.12')

The  $Lomax(a, \theta)$  distribution has the pdf in the form

$$f(x, a, \theta) = \begin{cases} 0 & \text{if } x \le 0\\ \frac{a\theta}{(1+\theta x)^{a+1}} & \text{if } x > 0. \end{cases}$$
(1.13)

The Lomax cdf is

$$F(x, a, \theta) = \begin{cases} 0 & \text{if } x \le 0\\ 1 - \frac{1}{(1+\theta x)^a} & \text{if } x > 0 \end{cases}$$
(1.14)

and its inverse (the solution of equation F(x) = U) is

$$F^{-1}(U) = \frac{1}{\theta} \left( \left(\frac{1}{U}\right)^{\frac{1}{a}} - 1 \right), \text{ for } 0 < U < 1.$$
 (1.14')

If U is an uniform random number over (0, 1), then the target distributions are simply simulated by the *inverse method* as

$$X = F^{-1}(U),$$

where  $F^{-1}$  is either given by (1.12') or by (1.14').

## 2. RESULTS: MIXED DISTRIBUTIONS OBTAINED

The pdf's of mixed distributions are in the form

$$f_V(x) = \sum_{k=1}^{\infty} P(N^* = k) k \varphi(x) (1 - \phi(x))^{k-1}, \qquad (2.1)$$

$$f_W(x) = \sum_{k=1}^{\infty} P(N^* = k) k \varphi(x) (\phi(x))^{k-1}.$$
 (2.2)

#### 2.1 The case of truncated $Poisson(\lambda)$ distribution

For the Poisson distribution Kus [2] derived a new distribution (of W), when the target distribution is *extreme value distribution*.In [1], Lupu derived distributions on V and W when the target distribution is exponential  $Exp(\lambda)$ .

Here, we assume first that the target distribution is  $Weibull(0, 1, \nu)$ ,  $\nu > 0$ , i.e a Weibull standard which has the pdf

$$\varphi(x) = \nu x^{\nu - 1} e^{-x^{\nu}}, x > 0.$$
(2.3)

Therefore, for the variable V we have

$$f_V^{(PW)}(x) = \sum_{k=1}^{\infty} \frac{1}{e^{\lambda} - 1} \frac{\lambda^k}{k!} k \nu x^{\nu - 1} e^{-x^{\nu}} e^{-(k-1)x^{\nu}} =$$
$$= \frac{\lambda}{e^{\lambda} - 1} \nu x^{\nu - 1} e^{-x^{\nu}} \sum_{k=1}^{\infty} \frac{(\lambda e^{-x^{\nu}})^{k-1}}{(k-1)!}$$

which, after some calculations, gives finally

$$f_V^{(PW)}(x) = \frac{\lambda}{e^{\lambda} - 1} \nu x^{\nu - 1} e^{-x^{\nu}} e^{\lambda e^{-x^{\nu}}}.$$
 (2.4)

(The upper script (PW) means Poisson - Weibull and the lower script V refers to the random variable in (1.1)).

One can see that the corresponding cdf is

$$F_V^{(PW)}(x) = \frac{1}{e^{\lambda} - 1} (e^{\lambda} - e^{\lambda e^{-x^{\nu}}}).$$
(2.5)

For the random variable W we have

$$f_W^{(PW)}(x) = \sum_{i=1}^{\infty} \frac{1}{e^{\lambda} - 1} \frac{\lambda^k}{k!} k\nu x^{\nu - 1} e^{-x^{\nu}} (1 - e^{-x^{\nu}})^{k - 1}$$

which after some calculation gives

$$f_W^{(PW)}(x) = \frac{\lambda}{e^{\lambda} - 1} \nu x^{\nu - 1} e^{-x^{\nu}} \sum_{i=1}^{\infty} \frac{\lambda (1 - e^{-x^{\nu}})^{k-1}}{(k-1)!}$$

i.e.

$$f_W^{(PW)}(x) = \frac{\lambda}{e^{\lambda} - 1} \nu x^{\nu - 1} e^{-x^{\nu}} e^{\lambda(1 - e^{-x^{\nu}})}.$$
 (2.6)

(Here again, the upper script (PW) refers to Poisson - Weibull mixture and the lower script W refers to the random variable W in (1.1)).

In this case the cdf is

$$F_W^{(PW)}(x) = \frac{e^{\lambda e^{-x^{\nu}}} - 1}{e^{\lambda} - 1}.$$
(2.7)

## 2.2 The case of a truncated Geometric(p) distribution

The calculations below use the following elementary result Lemma 1 If  $0<\rho<1$  then

$$\sum_{k=1}^{\infty} k \rho^{k-1} = \frac{1}{(1-\rho)^2}.$$

In this case we deal only with  $Weibull(0, 1, \nu)$  target distribution and with  $Pareto(a, \theta)(Lomax(a, \theta))$  target distribution.

• In the case of Geometric(p) distribution, the random variable V in (1.1) when the target distribution is Weibull has the pdf

$$f_V^{(GW)}(x) = \sum_{i=1}^{\infty} pq^{k-1}k\nu x^{\nu-1}e^{-x^{\nu}}(e^{-x^{\nu}})^{k-1}.$$

Using Lemma 1, after some calculation one obtain

$$f_V^{(GW)}(x) = p\nu x^{\nu-1} \frac{e^{-x^{\nu}}}{(1 - qe^{-x^{\nu}})^2}.$$
 (2.8)

(Here the meaning of upper script (GW) is obvious). The corresponding cdf is

$$F_V^{(GW)}(x) = \frac{1}{q} - \frac{p}{q} \cdot \frac{1}{1 - e^{-x^{\nu}}}.$$
(2.8')

The mixture of geometric distribution with random variable W when the target distribution is Weibull has the pdf

$$f_W^{(GW)}(x) = p \sum_{i=1}^{\infty} k\nu x^{\nu-1} e^{-x^{\nu}} q^{k-1} (1 - e^{-x^{\nu}})^{k-1}$$

which after some calculations based on Lemma 1, gives

$$f_W^{(GW)}(x) = p\nu x^{\nu-1} e^{-x^{\nu}} \frac{1}{(p+qe^{-x^{\nu}})^2}.$$
 (2.9)

After some calculations the corresponding cdf is

$$F_W^{(GW)}(x) = \frac{p}{q} \frac{1}{p + q e^{-x^{\nu}}} - \frac{p}{q}.$$
 (2.9')

• Let us consider now the  $Lomax(a, \theta)$  target distribution which has the pdf

$$f(x) = \frac{a\theta}{(1+\theta x)^{a+1}}, \ x > 0 \ a > 0, \ \theta > 0.$$
(2.10)

The mixture distribution of the Geometric(p) with the random variable V, when the target distribution is  $Lomax(a, \theta)$  has the pdf

$$f_V^{(GL)}(x) = \sum_{i=1}^{\infty} pq^{k-1} \cdot \frac{a\theta}{(1+\theta x)^{a+1}} \frac{k}{((1+\theta x)^a)^{k-1}}$$

Using Lemma 1, after saome calculation we obtain

$$f_V^{(GL)}(x) = \frac{pa\theta(1+\theta x)^{a-1}}{[(1+\theta x)^a - q]^2}, x > 0.$$
 (2.11)

The *cdf* of  $f_V^{(GL)}(x)$  is (after some calculations!)

$$F_V^{(GL)}(x) = 1 - \frac{p}{(1+\theta x)^a - q}, \ x > 0.$$
(2.12)

The mixture distribution of the Geometric(p) with the random variable W, when the target distribution is  $Lomax(a, \theta)$  has the pdf

$$f_W^{(GL)}(x) = \sum_{i=1}^{\infty} pq^{k-1} \frac{ka\theta}{(1+\theta x)^{a+1}} \left(1 - \frac{1}{(1+\theta x)^a}\right)^{k-1}, \ x > 0.$$

Using again Lemma 1, after some calculations, we obtain the final form

$$f_W^{GL}(x) = pa\theta \frac{(1+\theta x)^{a-1}}{(p(1+\theta x)^a + q)^2}, \ x > 0.$$
(2.13)

The corresponding cdf of  $f_W^{(GL)}(x)$  is

$$F_W^{(GL)}(x) = 1 - \frac{1}{p(1+\theta x)^a + q}, \ x > 0.$$
(2.14)

### 3. SIMULATION OF INTRODUCED DISTRIBUTIONS

#### 3.1 Direct simulation as composition algorithms

Since all distributions are *mixture* distributions, they can be simulated by *composition algorithms* in the form

**begin** Simulate n a sampling value of  $N^*$ ; Simulate  $L_1, L_2, ..., L_n$  having the target distribution; Calculate V (or W) as in (1.1); Take X = V (or X = W) end.

The random variable X produced by the algorithm has the corresponding mixture distribution. Implementation of all algorithms does not imply any difficulty.

## 3.2 Simulation based on the inverse method

For all introduced distributions the cdf's were specified. Therefore simulation of any random variable X which has the cdf F(x) is done according to the following *inverse* algorithm

#### begin

Simulate a random number U uniform over (0,1); Calculate  $X = F^{-1}(U)$ ; end.

(In the algorithm,  $F^{-1}$  is the inverse function of F).

In the following we list-up the inverse  $F^{-1}$  of cdf's introduced in section 2.

1. The inverse of the  $cdf \ F = F_V^{(PW)}$  is

$$F^{-1}(U) = \{ -\log[\log[e^{\lambda} - U(e^{\lambda} - 1)]^{\frac{1}{\lambda}}] \}^{\frac{1}{\nu}}.$$
 (3.1)

2. The inverse of the  $cdf \ F = F_W^{(PW)}$  is

$$F^{-1}(U) = \{-\log[\log[1 + U(e^{\lambda} - 1)]^{\frac{1}{\lambda}}]\}^{\frac{1}{\nu}}.$$
(3.2)

3. The inverse of the  $cdf \ F = F_V^{(GW)}$  is

$$F^{-1}(U) = \left[-\log\left(\frac{q(1-U)}{1-qU}\right)\right]^{\frac{1}{\nu}}.$$
(3.3)

4. The inverse of the  $cdf \ F = F_W^{(GW)}$  is

$$F^{-1}(U) = \left[ -\log(\frac{q(1-U)}{p+qU} \right]^{\frac{1}{\nu}}.$$
(3.4)

5. The inverse of the *cdf*  $F = F_V^{(GL)}$  is

$$F^{-1}(U) = \frac{1}{\theta} \left[ \left( \frac{1 - qU}{1 - u} \right)^{\frac{1}{a}} - 1 \right].$$
 (3.5)

6. The inverse of the  $cdf \ F = F_W^{(GL)}$  is

$$F^{-1}(U) = \frac{1}{\theta} \left[ \left( \frac{1 - qU}{1 - U} \right)^{\frac{1}{a}} - 1 \right].$$
 (3.6)

For implementation of these simulation procedures, it must take care to avoid calculation of the log function when its *argument* is close to zero (i.e. < 0.0000001!).

#### 3.3 Simulation based on rejection method

In this subsection we use the following known result [4,5]

**Theorem 1** Assume that X is a random variable to be simulated and its pdf is f(x). Assume that Y is another random variable for wich a simulation method is known and its pdf is h(x), such that functions f and g have the same support. Assume also that exists a finite constant  $\alpha > 0$ , such as

$$\frac{f(x)}{h(x)} \le \alpha. \tag{3.7}$$

Let U be an uniform (0,1) variate stochastically independent from Y. Then, the conditional pdf of Y given that

$$0 < U \le \frac{f(Y)}{\alpha h(y)},\tag{3.8}$$

is f. (Therefore X is simulated by Y which satisfies condition (3.8))

From this Theorem results the following general rejection algorithm for simulating the random variable X:

## Algorithm R

#### repeat

Generate U an uniform random number over (01);

Generate Y having pdf h, such as Y and U are stochastically independent;

**until**  $U \leq \frac{f(Y)}{\alpha h(Y)}$ ; (This is the condition C). Take X = Y. It is shown that the probability to come out from the cycle **repeat... until** (i.e the *acceptance probability*), is

$$p_a = \frac{1}{\alpha},\tag{3.9}$$

therefore  $1 < \alpha < \infty$  and the algorithm is fast if  $p_a$  is close to one.

In the following we apply this theorem for simulating, by the rejection method, the distributions introduced in section 2. In fact will be presented only important features necessary to implement **Algorithm R** for each distribution, namely:

- The enveloping density h(x);
- The constant  $\alpha$ ;
- The acceptance probability  $p_a$ ;
- The simplified expression of the condition C in the cycle **repeat... until** C.
  - A note regarding selection of the *enveloping density*:we will try (and it will work!) as *enveloping densities* the densities corresponding to *target distributions*.
    - a. The distribution  $F_V^{(PW)}$ . The enveloping density selected is the standard  $Weibull(0, 1, \nu)$  density

$$h(x) = \nu x^{\nu - 1} e^{-x^{\nu}}, \quad x > 0, \quad \nu > 0.$$
(3.10)

The enveloping constant is

$$\alpha = \frac{\lambda e^{\lambda}}{e^{\lambda} - 1} \tag{3.10'}$$

and the *acceptance probability* is

$$p_a = \frac{e^{\lambda} - 1}{\lambda e^{\lambda}}.$$
(3.10")

The simplified form of the predicate C is

$$C = true, iff, U \le e^{-\lambda + \lambda e^{-Y^{\nu}}}.$$
(3.10''')

b. The distribution  $F_W^{(PW)}$ . The enveloping density is the same as in **a**.; the constant  $\alpha$  and the acceptance probability are also the same. The final form of the predicate C is

$$C = true, iff, U \le e^{-\lambda + \lambda(1 - e^{-Y^{\nu}})}.$$
(3.11)

Even if the probability  $p_a$  is the same for cases **a**. and **b**., the rejection algorithm in the case **a**. is faster because the condition C in this case requires a smaller time-complexity.

c. The distribution  $F_V^{(GL)}$ . The enveloping density h(x) is also given by (3.10); the constant  $\alpha$  and the probability  $p_a$  are

$$\alpha = \frac{1}{p}, \quad p_a = p. \tag{3.12}$$

The final form of the predicate C is

$$C = true, \, iff \, U \le \frac{p^2}{(1 - qe^{-Y^{\nu}})^2}.$$
 (3.12')

d. The distribution  $F_W^{(GW)}$ . The enveloping density h(x) is also (3.10); The constants  $\alpha$  and  $p_a$  are the same as in case **c.** and the predicate C is

$$C = true, iff U \le \frac{p^2}{(p + qe^{-Y^{\nu}})^2}.$$
 (3.13).

Note that complexity of the two rejection algorithms (cases **c**. and **d**.) differs only by the time necessary to evaluate predicates C.

e. The distribution  $F_V^{(GL)}$ . The enveloping density selected is the  $Lomax(a, \theta)$  density, i.e.

$$h(x) = \frac{a\theta}{(1+\theta x)^{a+1}}, x > 0, a, \theta > 0.$$
(3.14)

The enveloping constant  $\alpha$  and the acceptance probability are

$$\alpha = \frac{1}{p}, \, p_a = p, \tag{3.14'}$$

and the final form of the predicate C is

$$C = true, \, iff \, U \le \frac{p^2 a \theta (1 + \theta Y)^{2a}}{\left[(1 + \theta Y)^a - q\right]^2}.$$
 (3.14")

f. he distribution  $F_W^{(GL)}$ . The enveloping density is also  $Lomax(a, \theta)$  and the constants  $\alpha, p_a$  are the same as in the case e., i.e. are in the form (3.14'). The final form of the predicate C is

$$C = true, \, iff \, U \le \frac{p^2 a \theta (1 + \theta Y)^{2a}}{[p(1 + \theta Y)^a + q]^2}.$$
(3.15)

If the Weibull and Lomax random variates are simulated by the inverse methods, then the final form of predicates could be further simplified, thus reducing execution time. • Note that the **moments** of mixed distributions introduced exist and they could be simply calculated.

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